

# Midterm I Exam Review

15-122 Principles of Imperative Computation  
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## 1 Modular Arithmetic

In C0, values of type `int` are defined to have 32 bits. In this problem we work with a version of C0 where values of type `int` are defined to have only 9 bits. In other respects it is the same as C0. All integer operations are still in two's complement arithmetic, but now modulo  $2^9$ . All bitwise operations are still bitwise, except on only 9 bit words instead of 32 bit words.

**Task 1.** Fill in the missing quantities, in the specified notation.

- a. The minimal negative integer, in decimal: \_\_\_\_\_
- b. The maximal positive integer, in decimal: \_\_\_\_\_
- c.  $-4$ , in hexadecimal: `0x` \_\_\_\_\_
- d.  $44$ , in hexadecimal: `0x` \_\_\_\_\_
- e. `0x49`, in decimal: \_\_\_\_\_

**Task 2.** Assume `int x` and `int y` have been declared and initialized to unknown values. For each of the following, indicate if the expression always evaluates to true, or if it could sometimes be false. In the latter case, indicate a counterexample in the C0 dialect described here by giving a value for `x` and `y` that falsifies the claim. You may use decimal or hexadecimal notation. You may use the functions `int_min()` and `int_max()` and assume they are correct for 9 bit ints.

- a. `x >= x - 1` \_\_\_\_\_
- b. `(~x) ^ (x) == -1` \_\_\_\_\_
- c. `x+(y+1)-2*(x-1)-3 == -x+y` \_\_\_\_\_
- d. `(x!=-x || y!==-y) || x==y` \_\_\_\_\_
- e. `((x<<1)>>1) | (x & 0x80) == x` \_\_\_\_\_
- f. `((x>>1)<<1) | (x&1) == x` \_\_\_\_\_
- g. `x <= (1<<(7-1))-1` \_\_\_\_\_
- h. `x+x == 2*x` \_\_\_\_\_

## 2 Ternary Search.

Consider the following variation of binary search algorithm. Instead of checking the middle element of the sorted array  $A[]$ , check the element at position  $n/3$ . Then proceed in the same way as in binary search. If you are looking for  $x$  then

- if  $x == A[n/3]$  you have found it.
- if  $x < A[n/3]$  you search the first third of the array, namely at indexes  $< n/3$ .
- if  $x > A[n/3]$  you search the rest two thirds of the array at indexes  $> n/3$ .

Consider the following implementation:

```
int tersearch(int x, int[] A, int n)
//@requires 0 <= n && n <= \length(A);
//@requires is_sorted(A, n);
/*@ensures (-1 == \result && !is_in(x, A, n))
    || ((0 <= \result && \result < n) && A[\result] == x);
@*/
{
    int lower = 0;
    int upper = n;
    while (lower < upper)
        //@loop_invariant 0 <= lower && lower <= upper && upper <= n;
        //@loop_invariant lower == 0 || A[lower-1] < x;
        //@loop_invariant upper == n || A[upper] > x;
        {
            int mid = lower + (upper-lower)/3;
            if (A[mid] < x)
                lower = mid+1;
            else if (A[mid] > x)
                upper = mid;
            else //@assert A[mid] == x;
                return mid;
        }
    //@assert lower == upper;
    return -1;
}
```

**Task 1.** Prove the correctness of this function, using the method presented in class.

**Task 2.** Find an expression for the worst-case runtime of ternary search. Compare this asymptotically to binary search.

### 3 Runtime Analysis

Determine the big O complexity of the following expression:

$$T(n) = 4n^5 + 10n^3 + \log n$$

Use the formal definition of big O to justify your answer.